

$\Sigma = \overbrace{\{a, b, c\}}^{\text{alphabet}}$   
symbols

string:  
sequence of symbols

Examples of strings over  $\Sigma$ :

abc a  $\epsilon$  b ccccc

~~abcd~~ not over  $\Sigma$   
because d is not in  $\Sigma$

A string is **over**  $\Sigma$  if all of its symbols are in  $\Sigma$

A **language** over  $\Sigma$  is a set of strings over  $\Sigma$

Examples of languages over  $\Sigma$ :

$\{aa, bb, cc\}$   $\{abc, a, \epsilon, b, ccccc\}$   $\{b^n \mid n \geq 0\}$   
 $b^2 = bb$   $b^3 = bbb$   
 $b^0 = \epsilon$

We can use regular expressions to describe these languages.

# Regular Expressions

Basis steps:

$a$  is a regular expression, for  $a \in \Sigma$

$\epsilon$  is a regular expression

$\emptyset$  is a regular expression

regular expressions

$$L(a) = \{a\}$$

string in the language

$$L(\epsilon) = \{\epsilon\}$$

$$L(\emptyset) = \{\emptyset\}$$

the empty set  $\rightarrow$  represent language with 0 strings in it

Recursive Steps:

$(R_1^*)$  is a regular expression when  $R_1$  is a regular expression

$(R_1 \circ R_2)$  is a regular expression when  $R_1$  and  $R_2$  are regular expressions

$(R_1 \cup R_2)$  is a regular expression when  $R_1$  and  $R_2$  are regular expressions

$$\Sigma = \{a, b, c\}$$

# Regular Expressions

$(R_1^*)$  is a regular expression when  $R_1$  is a regular expression

$$\begin{aligned} (b)^* & L(b) = \{b\} \\ & L(b^*) = \{b^n \mid n \geq 0\} \\ & \epsilon \quad b \quad bb \quad bbb \dots \end{aligned}$$

$$\begin{aligned} \Sigma^* & L(\Sigma^*) = \\ & \text{the set of strings} \\ & \text{over } \Sigma \end{aligned}$$

$(R_1 \circ R_2)$  is a regular expression when  $R_1$  and  $R_2$  are regular expressions

$$\begin{aligned} a \circ a \circ a \circ a \circ a & L(a \circ a \circ a \circ a \circ a) = \\ \text{(shorthand: } a \circ a \circ a \circ a \circ a) & \{a \circ a \circ a \circ a \circ a\} \end{aligned}$$

$$\begin{aligned} c \circ b^* & L(c \circ b^*) \\ \text{(shorthand: } c \circ b^*) & = \{c b^n \mid n \geq 0\} \\ & c \quad cb \quad cbb \quad cbbb \dots \end{aligned}$$

$(R_1 \cup R_2)$  is a regular expression when  $R_1$  and  $R_2$  are regular expressions

$$\begin{aligned} a \cup c & L(a \cup c) \\ & = L(a) \cup L(c) \\ & = \{a, c\} \end{aligned}$$

$$\begin{aligned} c b^* \cup a \circ a \circ a \circ a & L(c b^* \cup a \circ a \circ a \circ a) \\ & = L(c b^*) \cup L(a \circ a \circ a \circ a) \\ & = \{c b^n \mid n \geq 0\} \cup \{a \circ a \circ a \circ a\} \end{aligned}$$

$$\Sigma = \{a, b, c\}$$

# Regular Expressions

$$\underline{cb^* \cup aaaaa}$$

$$L(cb^* \cup aaaaa) = \{cb^n \mid n \geq 0\} \cup \{aaaaa\}$$

Implicitly:

cb cbbb bbaaaaa

$$(c(b^*)) \cup (aaaaa)$$

What happens when we evaluate in a different order?

$$\boxed{(cb)^*} \cup (aaaaa)$$

$$L((cb)^*) = \{(cb)^n \mid n \geq 0\}$$

cbcbcbcb  $\epsilon$

examples of strings

→ The language is not the same!

Precedence order: First \*, then  $\cup$ , then  $\cup$